

Internet clouds are (also) unpredictable!

A study of **flow rerouting** and **latency variations** across the largest Cloud network worldwide

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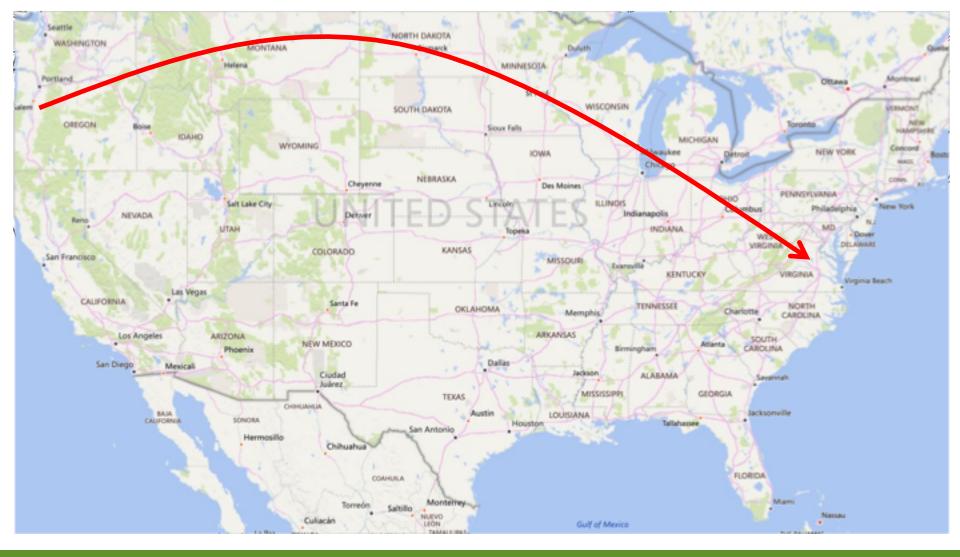
Joint work with: Waleed Reda, Kirill Bogdanov, Alexandros Milolidakis, Gerald Q Maguire Jr, Dejan Kostic



Thanks to the RACI initiative!



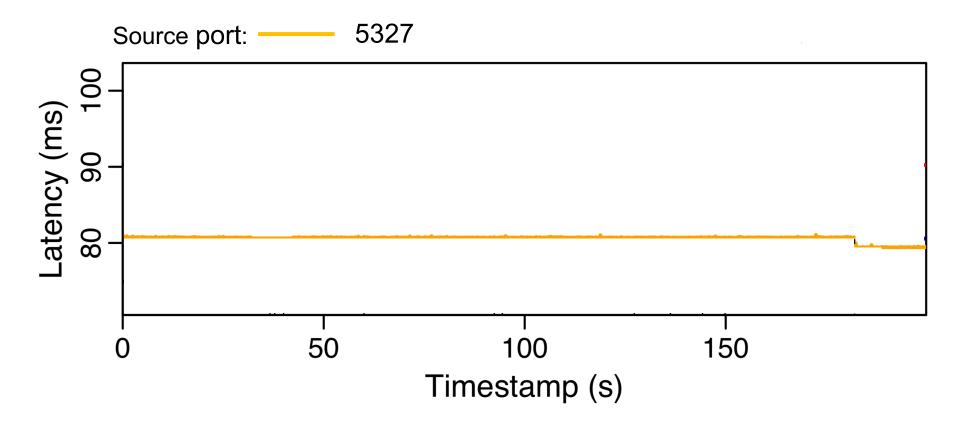
Let's open a TCP connection between Oregon and Virginia



* map of USA taken from Bing Maps

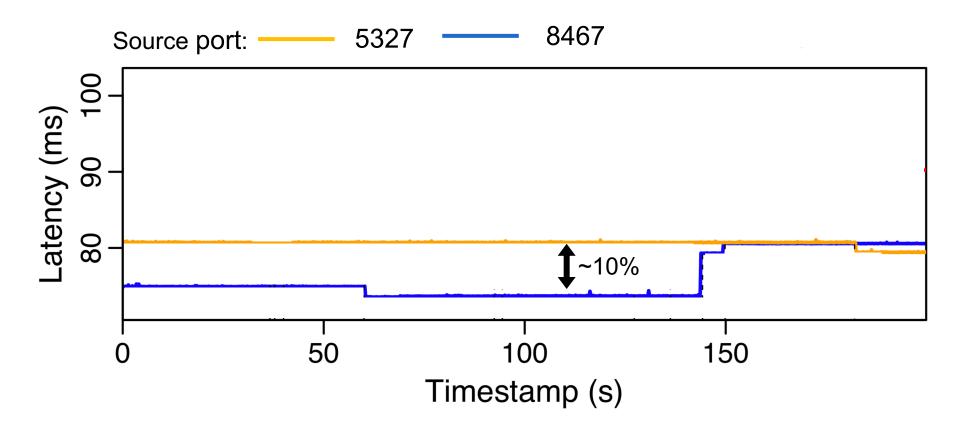


Stable Round Trip Time latency of roughly 80ms



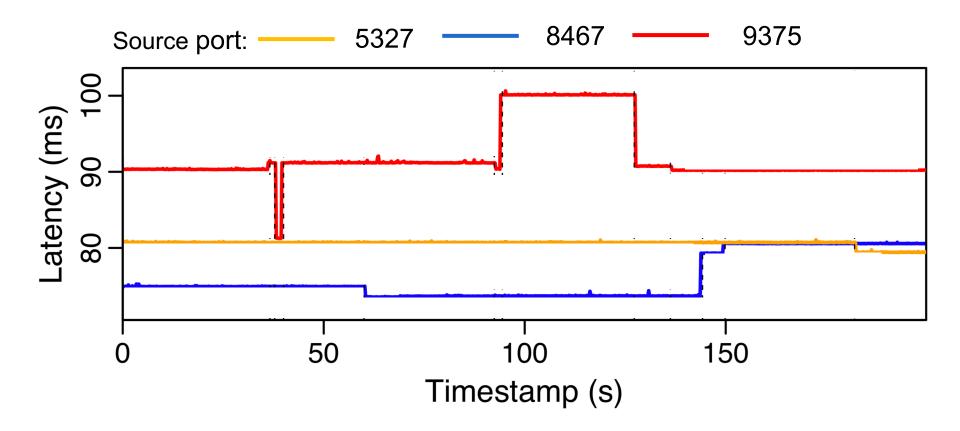


We open a second connection: Up to 10% better latency!



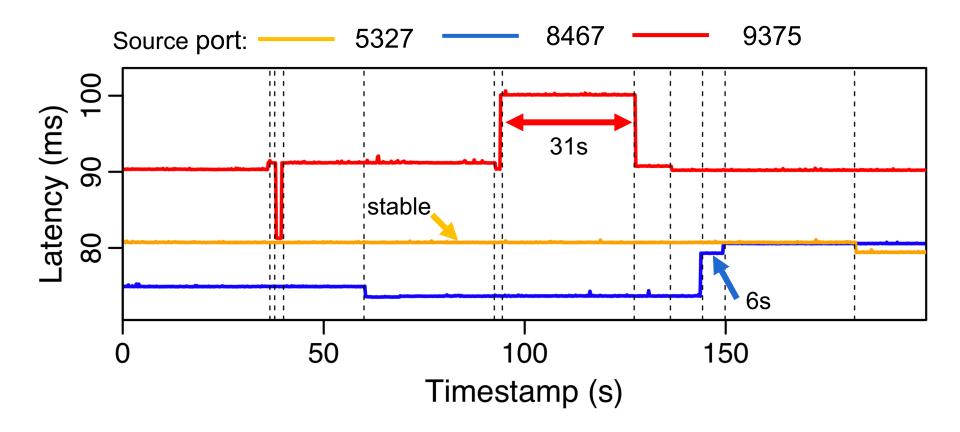


We open a third connection: Unstable and worse latency!



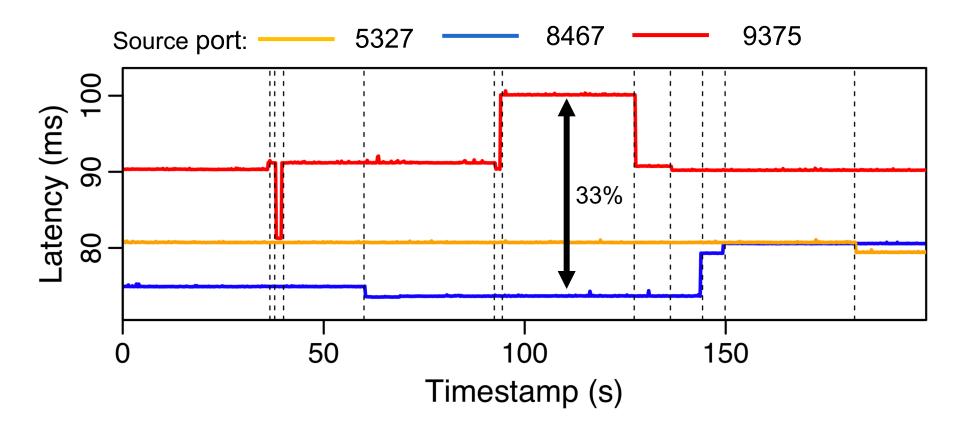


Insight #1: Flows experiences path changes at a second-level time scales



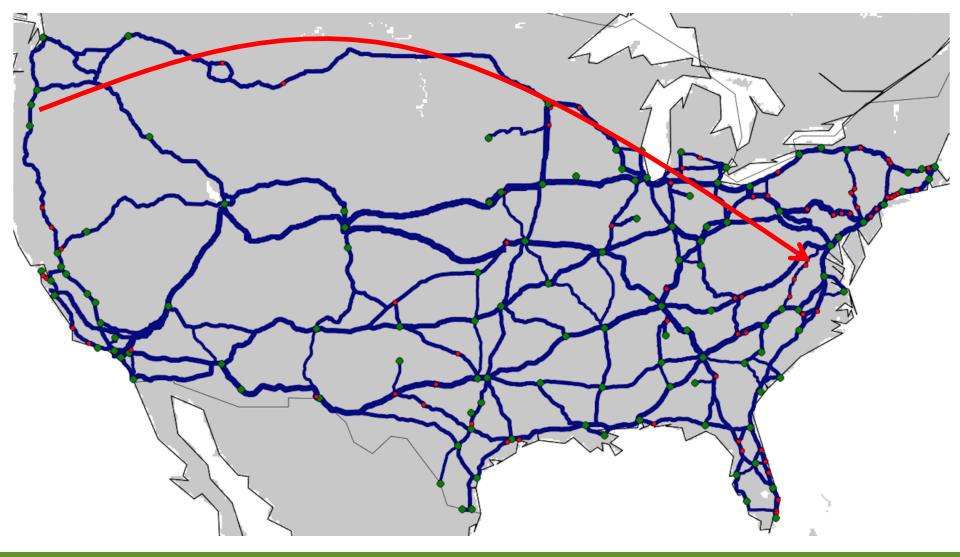


Insight #2: Flows experience unfair performance





Why do we see so many different latency behaviours?





Traffic Engineering (TE) : balance flows of traffic in a network

Periodic TE re-optimization in cloud networks:

- 2011: MPLS autobandwidth in Microsoft WAN*
- today: (claimed) **SDN-based** TE at Amazon, Google, Microsoft

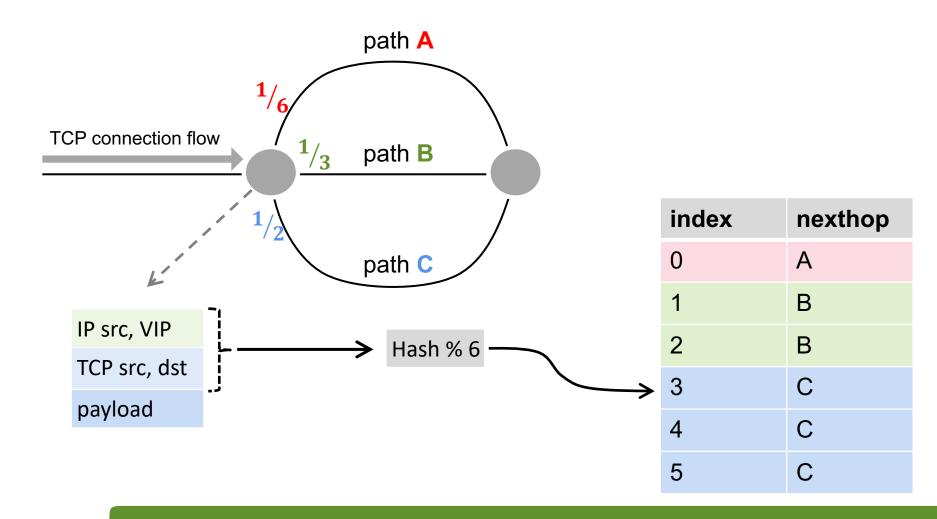
Two-phase TE:

- **compute** (multiple) **tunnels** among region pairs
- **compute** traffic **splitting ratios** among tunnels
 - typically hash-based implementation -> flow consistency
 - risk of **reshuffling flow packets** upon **splitting ratio update**!

* A. Pathak et al "Latency Inflation with MPLS-based Traffic Engineering" IMC 2011

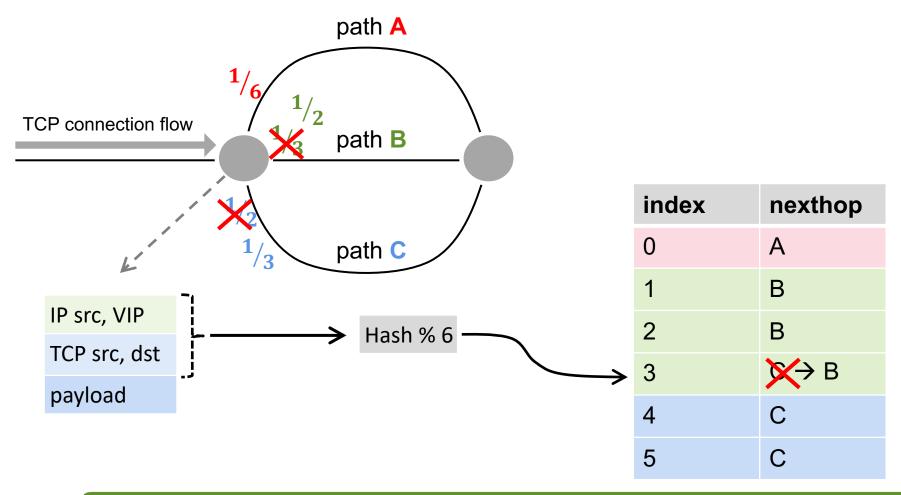


Stateless hash-based traffic splitting



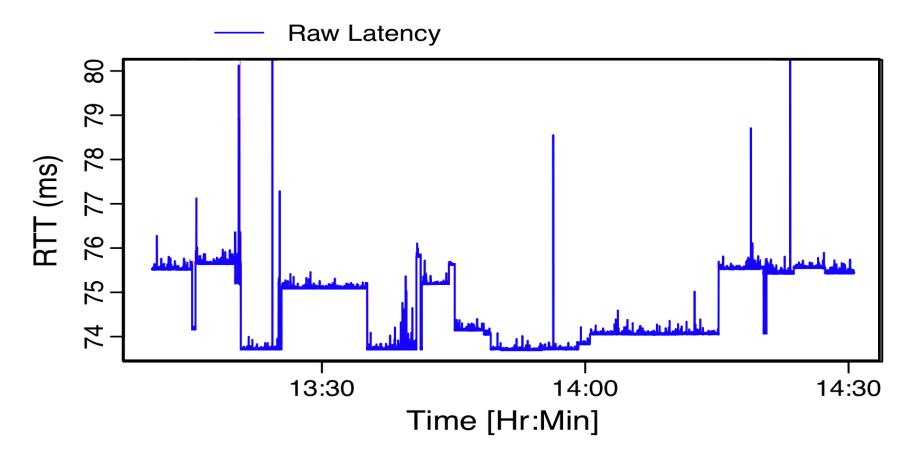


Stateless hash-based traffic splitting: packets belonging to the same connection may be rerouted when traffic splitting ratios are modified





Let's open a new connection and dissect it: latency = base propagation latency + congestion

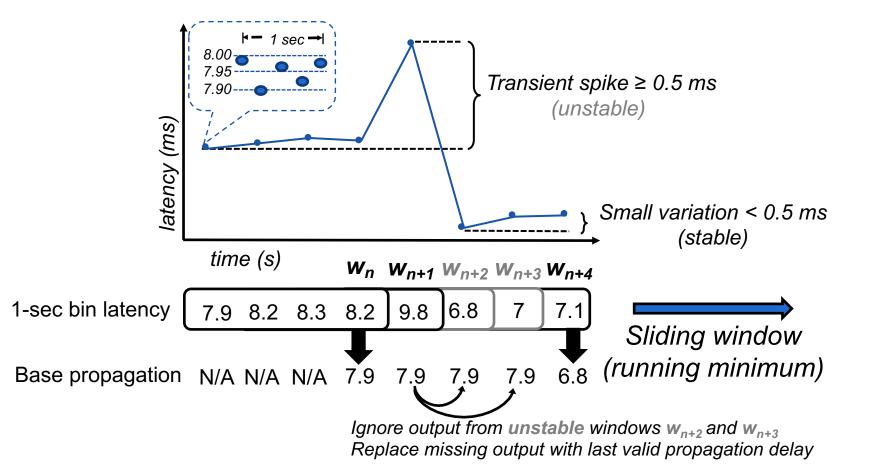




The technical yet necessary part (bear with me): Extracting base propagation latency

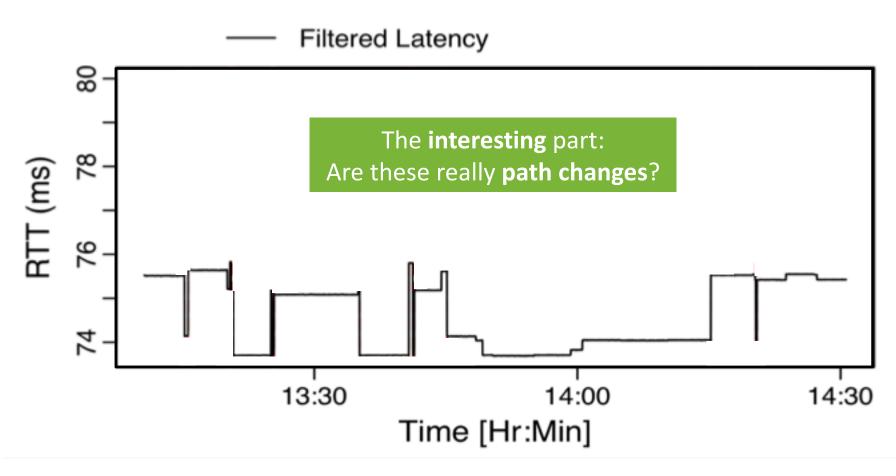
Rolling-minimum sliding window:

- stable if 3 out of 4 samples are within 0.5ms
- update base propagation latency only if stable AND >0.5ms change



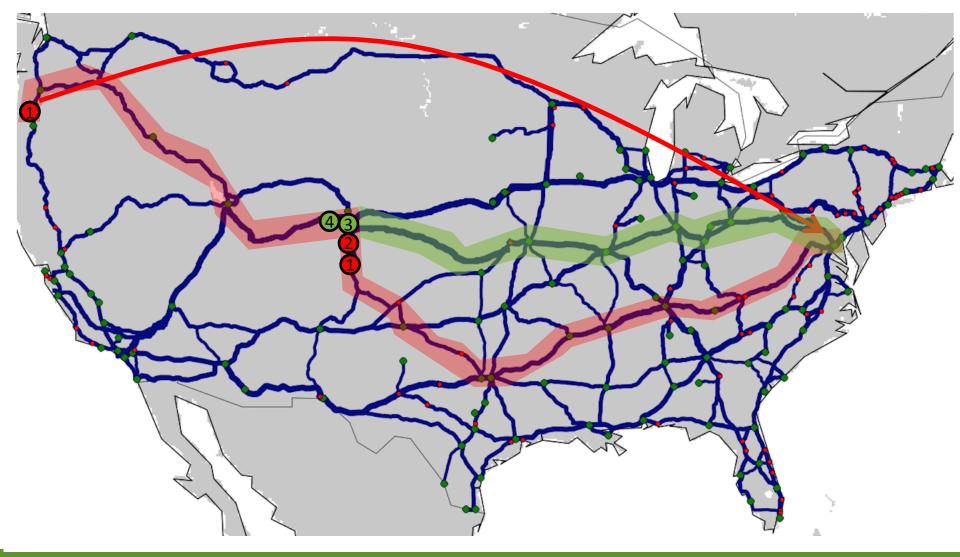


The filtered base propagation latency



KTH vetenskap och konst

Key intuition: a path change from a high to a low latency path causes packet reordering

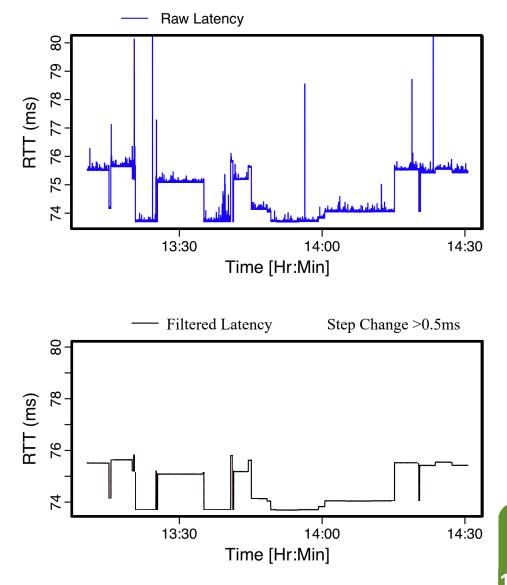




The path change detector is accurate and conservative

We **correlate** packet **reordering** events with propagation **latency**:

- zero false positives, i.e., a latency decrease >0.5ms is a path change
- limited false negatives, *i.e.*, path changes within 0.5ms
 -> conservative approach
- henceforth, latency means base propagation latency





We can now detect path changes

We set up an extensive measurement study of AWS

Config.	Macro-scale	Micro-scale
# of DC Pairs	120	4
# of Flows	512	512
Probing Rate	10 probes every 30s	5 probes/s
Flow Generation	Dynamic (every 30s)	Static
Duration	2 days	1 week
Ping Mechanism	Raw Sockets	TCP Ping



Picture credits: AWS re:Invent 2016: Tuesday Night Live with James Hamilton



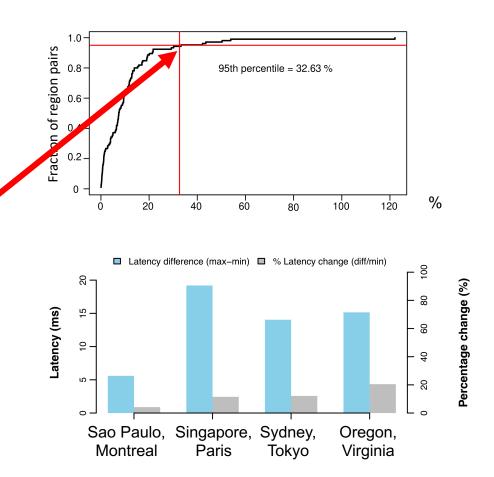
Some AWS regions have high latency variations across different flows (i.e., different source ports)

We looked at all the 120 region pairs in AWS:

- Latency inflation: max-latency/min-latency
- >32% latency inflation for 5% of region pairs

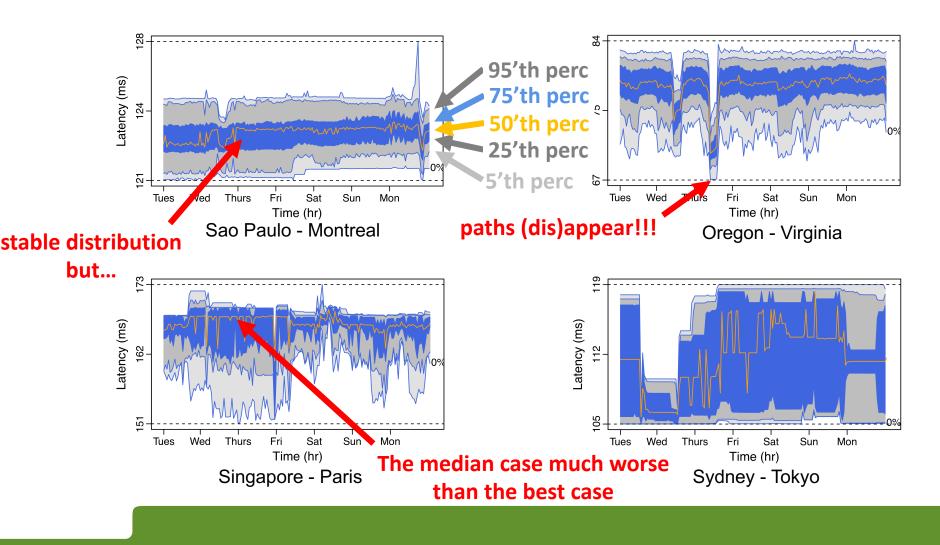
We zoomed into four AWS region pairs for the micro-scale experiments

 neither the worst nor the best pairs





The flow-latency spectrum over time for four different pairs of regions

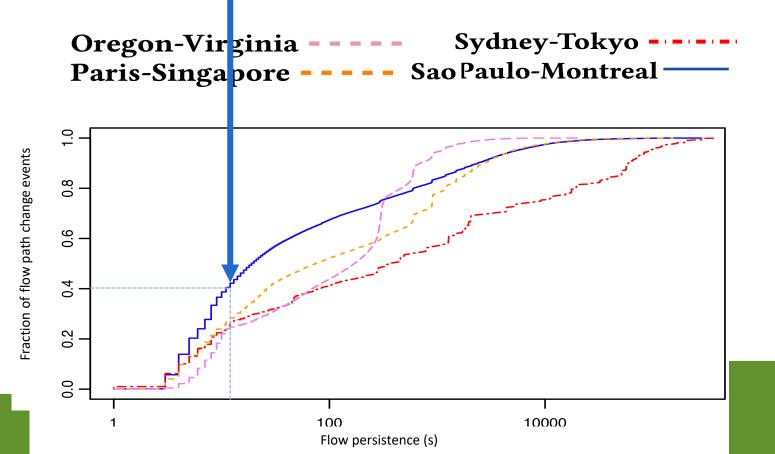




Stable distribution but... are the flows routed always on the same path? Not necessarily!

Consider all events when a flow changes path:

• 40% of the cases, the flow moves away within 10 seconds

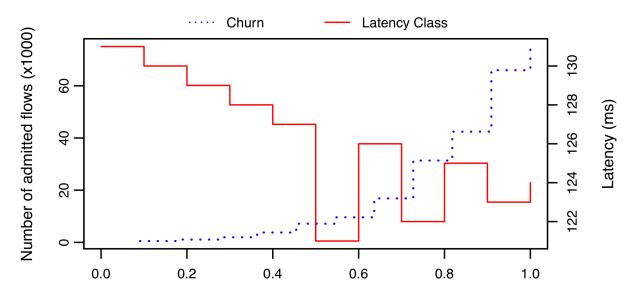




Flows routes on low-latency paths are more likely to experience a path change

We counted the number of flows moved away from a path per flow latency class (Oregon-Virginia):

- number of moved flows is **inversely proportional** to flow latency
- likely constrained **shortest-path-based TE**



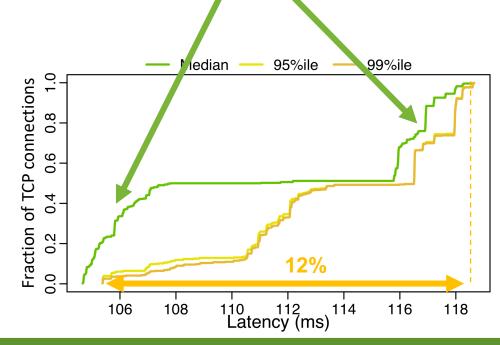
Cumulative fraction of latency classes



Measuring the level of "unfairness" among flows

We measure the flow latency of each flow in 20-sec interval bins

- we plot the 50'th, 95'th, 99'th latency percentiles
- bimodal behaviour at 50'th
- 12% difference for all percentiles

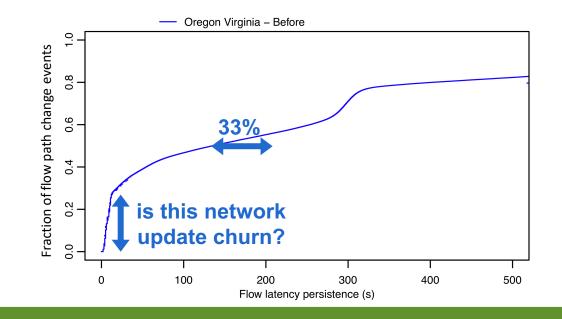




Measuring instabilities during "low" or "high" season

Large cloud networks tune the "aggressiveness" of their TE mechanisms according to the users' load We performed again the same measurements during "low"

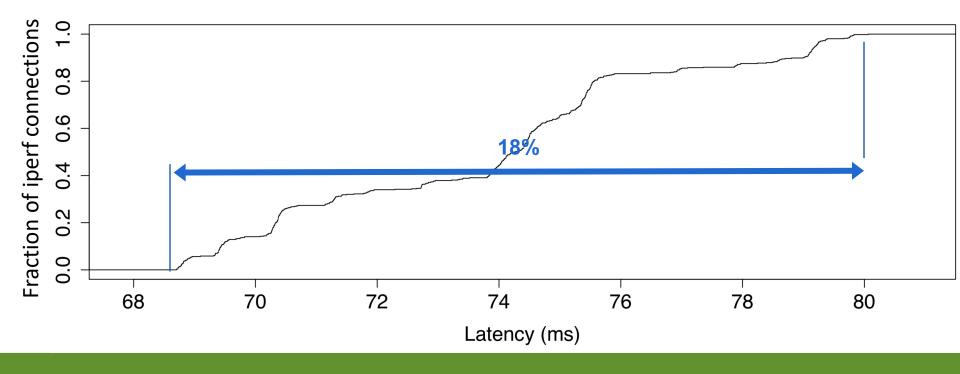
season, expecting to see less path changes





Moving 1K of data can take 18% more time when the flow is mapped to a high latency path

We generated a large number of iperf tests from Oregon (US west coast) to Virginia (US east coast) We moved 1K of data in each iperf test We did not observe any packet retransmission





Conclusions and takeaway

We measured the largest worldwide cloud backbone network (AWS):

- Insight #1: very reactive TE, especially across some regions
- Insight #2: flows experience unfair treatment
 - both in Round Trip Time and in the number of path changes

Far-reaching implications on the cloud $\leftarrow \rightarrow$ tenants interaction:

- congestion control: TCP Cubic suffers from packet reordering
- latency: low-latency geo-distributed emerging applications require low and deterministic latencies
- **selfish-routing**: application developers can force multiplexing of traffic on shortest paths, ultimately defeating cloud TE operation

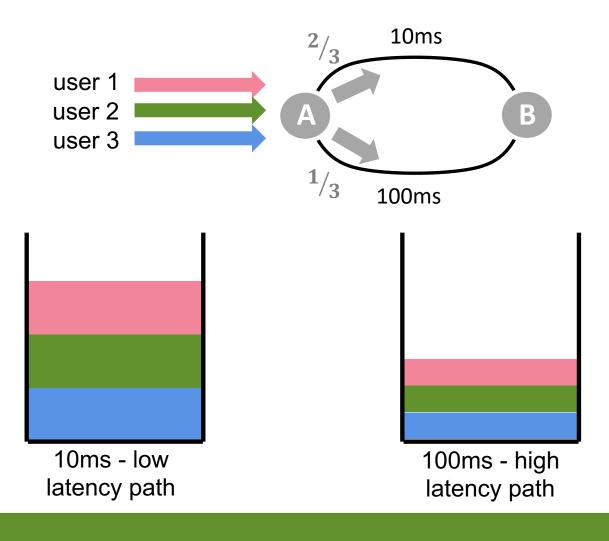
Thank you!

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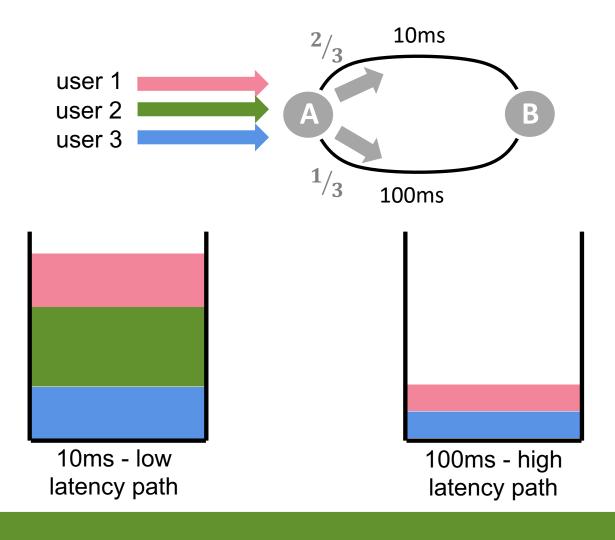


Selfish routing: giving control to the cloud tenants may lead to an unhealthy situation



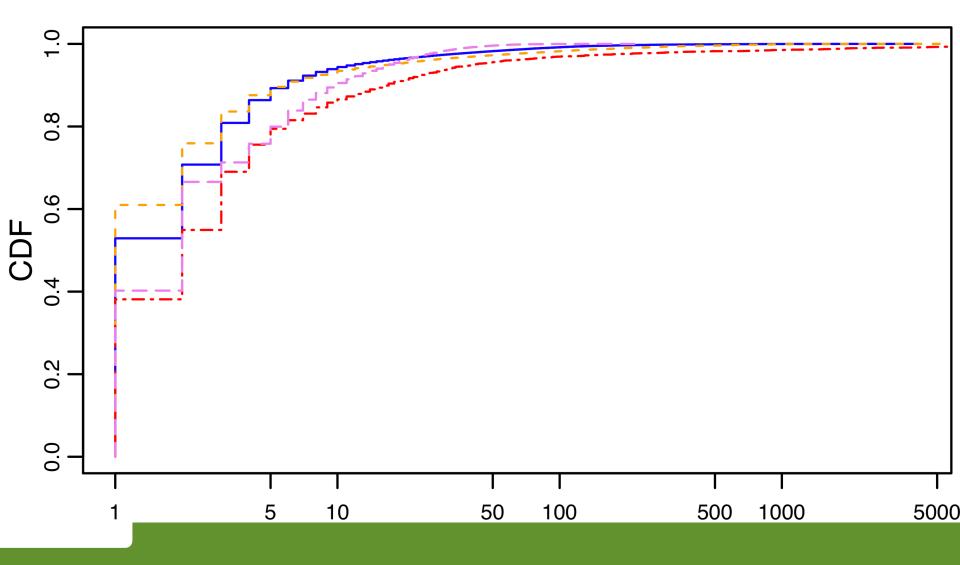


Selfish routing: giving control to the cloud tenants may lead to an unhealthy situation



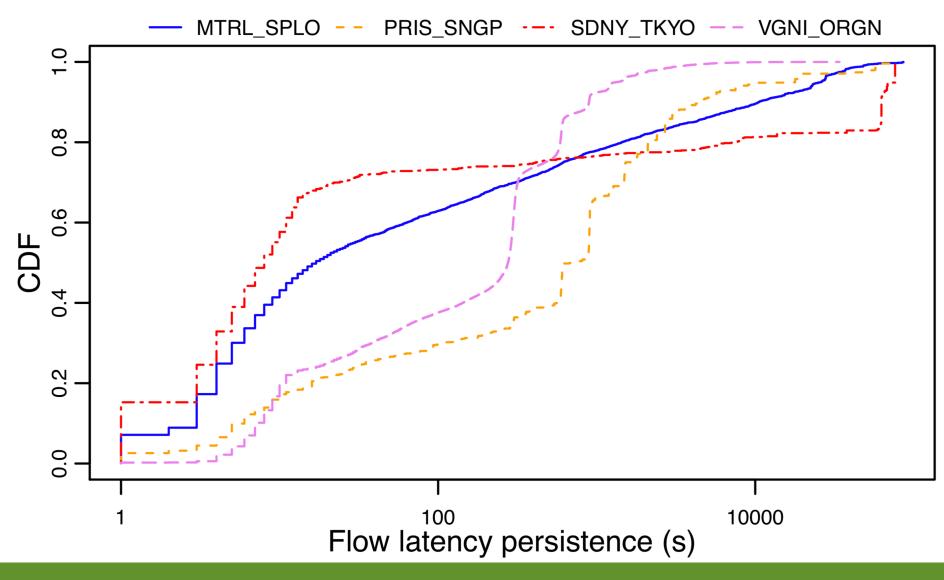


Flow path change interarrival time





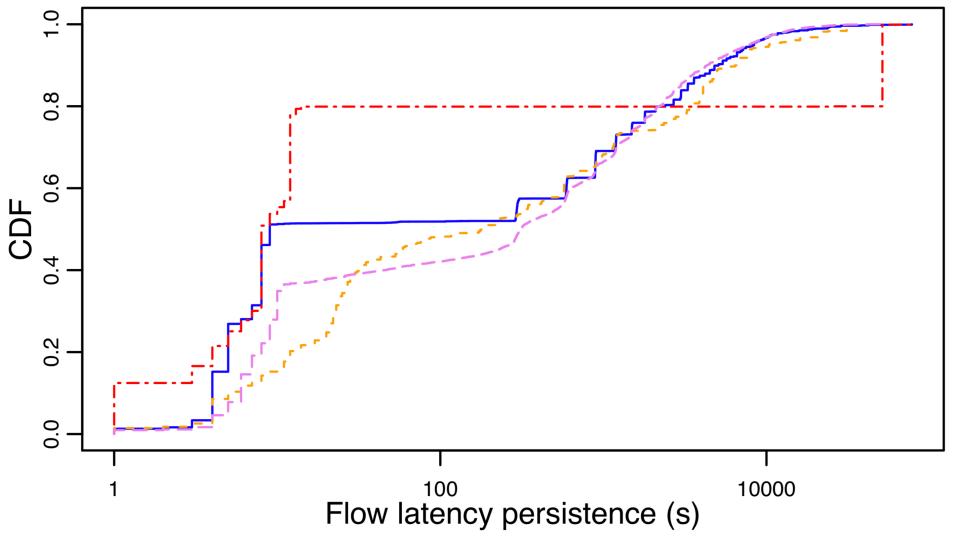
Flow persistence – forward path





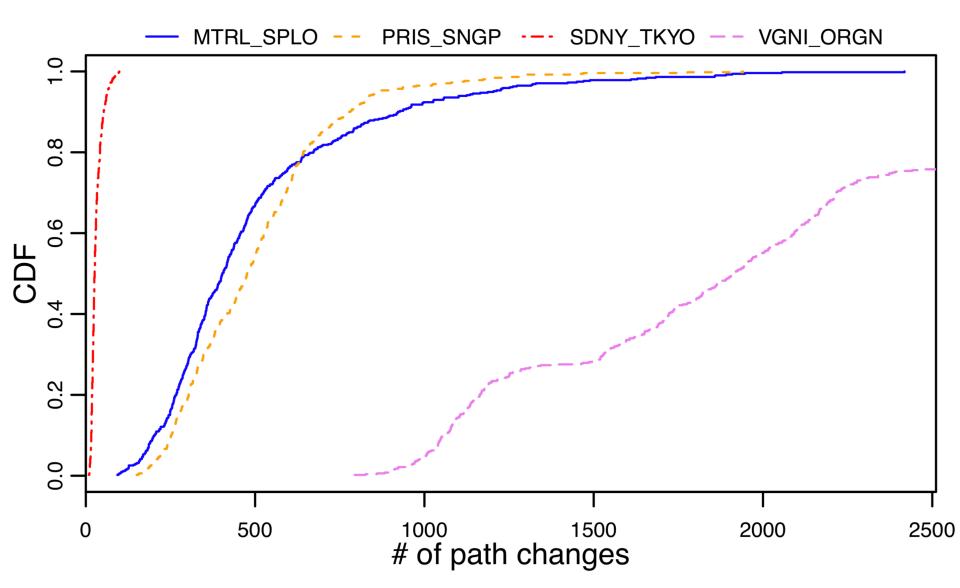
Flow persistence –reverse path

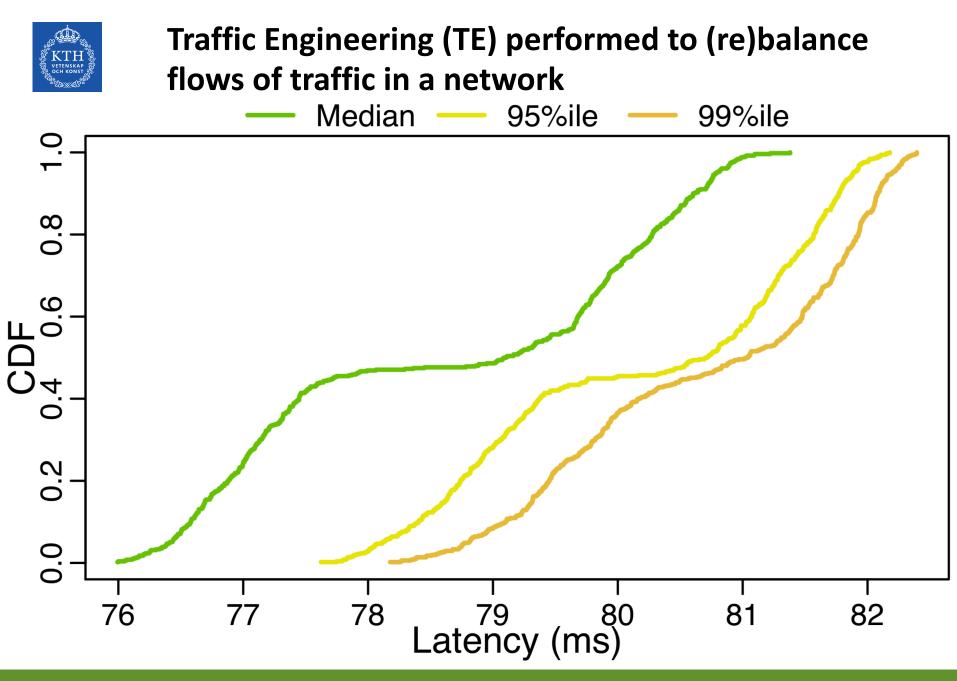
---- MTRL_SPLO - - PRIS_SNGP --- SDNY_TKYO -- VGNI_ORGN





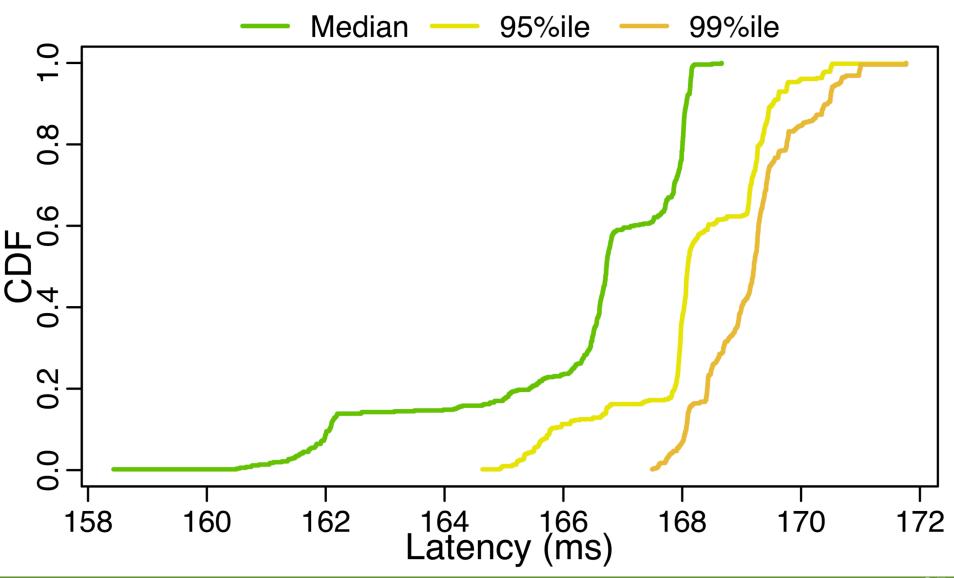
Traffic Engineering (TE) performed to (re)balance flows of traffic in a network

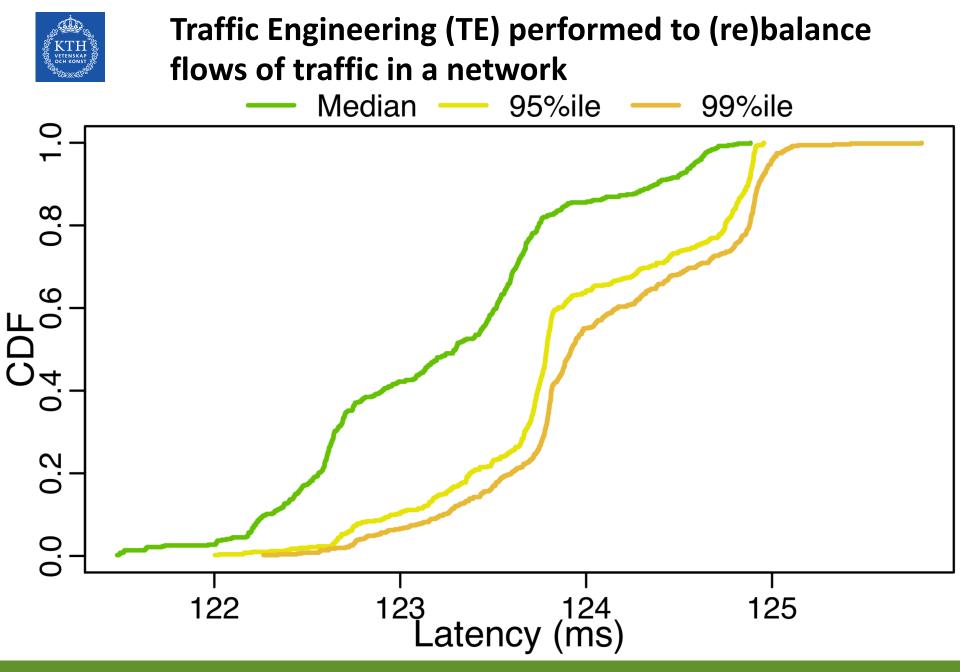






Traffic Engineering (TE) performed to (re)balance flows of traffic in a network





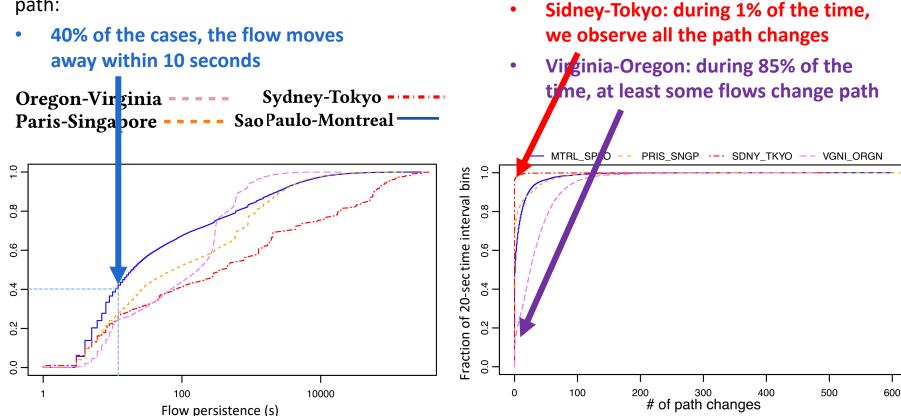


Fraction of flow path change events

Stable distribution but... are the flows routed always on the same path? Not necessarily!

Consider all events when a flow changes path:

We count #path-changes each 20 seconds





This talk in a nutshell

We measured the largest worldwide cloud backbone network (AWS):

- Insight #1: very reactive TE performed across datacenters
- **Insight #2**: tenants' flows experience **unfair** treatment
 - both in Round Trip Times and number of path changes

Far-reaching implications on how clouds and tenants interact:

- **congestion control:** TCP Cubic suffers from packet reordering
- latency: low-latency geo-distributed emerging applications require low and deterministic latencies
- **selfish-routing**: application developers can force multiplexing of traffic on low-latency paths, ultimately hindering cloud TE operation